# On the Complexity of Learning from Label Proportions

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#### Reference

# Benjamin Fish, Lev Reyzin. On the Complexity of Learning from Label Proportions. IJCAI 2017

#### A Motivating Example: Elections

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# Elections

- We know who voted.
- We know the result.
- But we don't know who voted for whom!
- We want to train a model that predicts future elections.

# Supervised Learning (PAC)

A class of functions H is **PAC learnable** if there is an efficient algorithm A such that for every target function c in H, any distribution D over  $\{0,1\}^n$ , and for any  $\varepsilon$ ,  $\delta > 0$ , given

 $m \ge poly(1/\epsilon, 1/\delta, n, size(c))$ 

labeled examples drawn i.i.d. from D, returns a hypothesis h in H such that

$$\mathsf{P}[\mathsf{1}_{\mathsf{C}(\mathsf{X})\neq\mathsf{h}(\mathsf{X})} \leq \varepsilon] \geq 1 - \delta.$$

#### Learning from Proportions (LLP)

A class of functions H is **PAC learnable from label proportions** if there is an efficient algorithm A such that for every target function c in H, any distribution D over  $\{0,1\}^n$ , and for any  $\varepsilon$ ,  $\delta > 0$ , given

#### $m \ge poly(1/\epsilon, 1/\delta, n, size(c))$

examples drawn i.i.d. from D and  $\hat{p}$ , returns a hypothesis h in H such that

 $\mathsf{P}[|\mathsf{p}(c)-\mathsf{p}(h)| \le \varepsilon] \ge 1-\delta.$ 

# Label Proportions vs PAC

A class of functions H is PAC learnable if there is an efficient algorithm A such that for every target function c in

H, any distribution D over  $\{0, 1\}^n$ , and for any  $\varepsilon, \delta > 0$ , given  $m \ge poly(1/\varepsilon, 1/\delta, n, size(c))$  examples drawn i.i.d. from D and their labels, returns a hypothesis h in H such that  $P[1_{C(x) \ne h(x)} \le \varepsilon] \ge 1 - \delta$ .

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every target function c in H, any distribution D over  $\{0, 1\}^n$ , and for any  $\varepsilon$ ,  $\delta > 0$ , given  $m \ge poly(1/\varepsilon, 1/\delta, n, size(c))$ examples drawn i.i.d. from D and  $\hat{p}$ , returns a hypothesis h in H such that  $P[|p(c) - p(h)| \le \varepsilon] \ge 1 - \delta$ .

# Main Question

What is the complexity of Learning from Label Proportions? And how does LLP learning relate to normal PAC learning?

#### An Occam's Razor Bound for LLP

For target function c, w.p. at least  $1 - \delta$ , for all  $h \in H$ ,

 $|p_{C} - p_{h}| \le |\hat{p}_{C} - \hat{p}_{h}| + \tilde{O}(\log(1/\delta) (VC(H)/m)^{1/2})$ 

## Results

- LLP  $\subsetneq$  PAC<sup>1</sup>
- VC-dimension hardness for LLP<sup>1</sup> (not a complete characterization)
- A nontrivial problem in LLP

under standard complexity-theoretic assumptions

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*proof.* involves a reduction from PAC to LLP. Idea is to make LLP distribution that forces all (and only the) + examples to be labeled + by LLP learner if its threshold is to be met.

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$$D'(x) = \begin{cases} \frac{m}{km+m-k} & \text{if } x \in S \text{ and } c(x) = 1\\ \frac{1}{km+m-k} & \text{if } x \in S \text{ and } c(x) = 0\\ 0 & \text{otherwise} \end{cases}$$

where k = number of positive examples and  $\varepsilon = 1/(2m^2)$ 

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reminder: the VC-dim is the maximum number of points a class of functions can shatter.

<u>Theorem</u>: Let C be a hypothesis class s.t. VC(C)  $\ge n^{\gamma}$  for some constant  $\gamma > 0$ . There is no efficient algorithm for PAC learning C from label proportions unless NP = RP.

*proof idea*: can reduce from subset sum to LLP learning any class with large VC dimension.

How? choose a set of shattered points, make LLP learner solve subset sum to get the "right" threshold.

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# An LLP-Learnable Class

# $\{h : \{1, \dots, 2^n\} \to \{0, 1\} : \max \\ h(i)=h(j)=1 ||i-j| \le k \}$



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For k=log(n), VC-dimension is super-constant, yet there is a polytime algorithm.

#### LLP is Also Easy for Nice Distributions



# Conclusions

- I presented the beginnings of a learning theory for LLP.
- LLP is the simplest case of multi-bag learning.
- Extensions include to multiclass learning and regression.
- Would also be interesting to develop practical LLP algorithms.